

# Labwork 3

March 8, 2021

**Goal:** to practice recursive thinking and the programming principles taught in Lecture 3.

## 1 Recap

Functions with one or more arguments of a recursive type should be defined by recursion on that argument.

- For list argument(s) there should be a base case when the list is empty, and a recursive case when the list is of the form `(cons v lst)`.
  - the recursive calls should take “simpler” argument values, to ensure termination of computation.
- For numeric argument(s), there should be
  - one or more base cases for small numeric values
  - one or more recursive cases which take simpler (smaller) numeric values.

### Exercise set 1

1. Define the function `(joinLists lst1 lst2)` which returns the results of joining lists `lst1` and `lst2`.
  - Example: `(joinLists '(1 2) '(a b))` should return `'(1 2 a b)`.
  - **Suggestion:** implement recursion on the length of `lst1`.
  - Side remark: Racket has the predefined function `(append lst1 ... lstn)` to join an arbitrary number of lists.
2. (LW3 from labwork 2) Define the numeric function `(digit-sum N)` which computes the sum of digits of the decimal representation of `N`. For example, `(digit-sum 726)` should return 15 because  $7 + 2 + 6 = 15$ . Racket has the following predefined functions for `a, b ∈ ℕ`:
  - `(quotient a b)` ; returns the quotient of dividing `a` by `b`
  - `(remainder a b)` ; returns the remainder of dividing `a` by `b`

3. (LW4 from labwork 2) Define `(flatten s1)` which takes as input a nested list of symbols, and returns the list of symbols contained in `s1` in the order in which they occur when `s1` is printed. Intuitively, `flatten` removes all the inner parentheses from its argument. For example:

```
> (flatten '(a b c))
'(a b c)
> (flatten '((a b) c (((d) e)))
'(a b c d e)
> (flatten '((a) () (b ()) () (c)))
'(a b c)
```

Note that nested lists of symbols are defined by the grammar

```
SL ::= null | (cons s SL) | (cons SL SL)
```

where `s` is a symbol.

## 2 Recursion by solving a more general problem

Sometimes, the best way to solve a problem is to find a solution to a more general problem, and use that solution to solve the original problem as a special case.

### Example

Suppose `von` is a vector of numbers. Define `(vector-sum von)` which returns the sum of elements of `von`, using the functions

- `(vector-length von)`: returns the length (number of elements) of `von`
- `(vector-ref von i)`: returns the `i`-th element of `von`; the elements of `von` are indexed starting from 0.

Instead of defining `(vector-sum von)`, we can define the more general function

```
(partial-vector-sum von n)
```

which computes the sum of elements with indexes from 0 to `n` in `von`.

Note that `(vector-sum von)` coincides with

```
(partial-vector-sum von (vector-length von))
```

```
(define (vector-sum von)
  (define (partial-vector-sum von n)
    (if (= n 0)
        0 ; nothing to add
        (+ (partial-vector-sum von (- n 1))
           (vector-ref von (- n 1)))))
  (partial-vector-sum von (vector-length von)))
```

Note that

1. `partial-vector-sum` is tail recursive, therefore it's computation is very fast.
2. The argument `von` of `partial-vector-sum` does not change its value. Since the function `partial-vector-sum` is defined in the body of `vector-sum`, we can optimize the definition of `partial-vector-sum` by eliminating argument `von`:

```
(define (vector-sum von)
  (define (partial-vector-sum n)
    (if (= n 0)
        0 ; nothing to add
        (+ (partial-vector-sum (- n 1))
           (vector-ref von (- n 1)))))
  (partial-vector-sum (vector-length von)))
```

## Exercise set 2

1. Let `(rev2 lst1 lst2)` be the function which returns the result of joining the reverse of list `lst1` with list `lst2`. For example, `(rev2 '(3 4) '(2 1))` yields `'(4 3 2 1)`. Note that, if `lst1` is not empty, then `(rev2 lst1 lst2)` returns the same result as

```
(rev2 (cdr lst1) (cons (car lst1) lst2))
```

- (a) Write a tail recursive definition of `rev2`.
  - (b) Define the function `(revList lst)` which computes the reverse of list `lst` as a special case of using the function `rev2`.
2. Let `(f2 a b c)` the function which computes  $c \cdot a^b$  when  $a, b, c$  are non-negative integers. Note that

$$c \cdot a^b = \begin{cases} c & \text{if } b = 0, \\ c \cdot (a^2)^{b/2} & \text{if } b > 0 \text{ is even,} \\ (c \cdot a) \cdot (a^2)^{(b-1)/2} & \text{if } b \text{ is odd.} \end{cases}$$

- (a) Write a tail recursive definition of `f2`.
  - (b) Define the function `(power a b)` which computes  $a^b$  for  $a, b \in \mathbb{N}$  as a special case of using the function `f2`.
3. Newton discovered the following method to compute  $\sqrt{a}$  when  $a$  is a non-negative number:  $\sqrt{a}$  is the limit of the sequence of numbers  $(x_n)_{n \in \mathbb{N}}$  where

$$x_0 = 1.0, \quad x_{n+1} = (x_n + a/x_n)/2 \quad \text{for all } n \in \mathbb{N}.$$

- (a) Define the function `(improve xn a)` which takes as inputs the values of  $x_n$  and  $a$  and returns the value of  $x_{n+1}$ .

- (b)  $x$  a *good enough* approximation of  $\sqrt{a}$  if  $|x^2 - a| \leq 0.000001$ . Define the boolean function (`good? x a`) which returns `#t` if the value of  $x$  is a good enough approximation of  $\sqrt{a}$ , and `#f` otherwise.
- (c) Newton's method finds a good approximation of  $\sqrt{a}$  starting from a number  $x$ , as follows: If  $x$  is a good approximation of  $\sqrt{a}$  it returns  $x$ , otherwise it returns a good approximation of  $\sqrt{a}$  starting from (`improve x a`).

Write a tail recursive definition of the function (`newton2 a x`) which uses Newton's method to compute a good approximation of  $\sqrt{a}$  starting from  $x$ .

4. Newton discovered the following method to compute  $\sqrt[3]{a}$  when  $a \in \mathbb{R}$ :  $\sqrt[3]{a}$  is the limit of the sequence of numbers  $(x_n)_{n \in \mathbb{N}}$  where

$$x_0 = 1.0, \quad x_{n+1} = (2 \cdot x_n + a/x_n^2)/3 \quad \text{for all } n \in \mathbb{N}.$$

- (a) Define the function (`improve3 xn a`) which takes as inputs the values of  $x_n$  and  $a$  and returns the value of  $x_{n+1}$ .
- (b)  $x$  a *good enough* approximation of  $\sqrt[3]{a}$  if  $|x^3 - a| \leq 0.000001$ . Define the boolean function (`good3? x a`) which returns `#t` if the value of  $x$  is a good enough approximation of  $\sqrt[3]{a}$ , and `#f` otherwise.
- (c) Newton's method finds a good approximation of  $\sqrt[3]{a}$  starting from a number  $x$ , as follows: If  $x$  is a good approximation of  $\sqrt[3]{a}$  it returns  $x$ , otherwise it returns a good approximation of  $\sqrt[3]{a}$  starting from (`improve3 x a`).

Write a tail recursive definition of the function (`newton3 a x`) which uses Newton's method to compute a good approximation of  $\sqrt[3]{a}$  starting from  $x$ .

5. Write a tail-recursive definition of the function (`eval v lst`) which takes as input a list of numbers

`lst = '(a0 a1 ... an)`

and computes the value of  $a_0 + a_1 \cdot v + \dots + a_n \cdot v^n$ .

6. Consider lists of symbols defined by the grammar

`SL ::= null | (cons s SL) | (cons SL SL)`

where  $s$  is a symbol. Write a tail recursive definition of the function (`flattenTR sl`) which computes the reverse of the flattened form of a list of symbols `sl`.

7. A rational number is a number whose value coincides with  $\frac{a}{b}$  where  $a, b \in \mathbb{Z}$  and  $b \neq 0$ . Suppose we choose to represent every rational number  $\frac{a}{b}$  as a pair (`cons a b`) where  $a, b \in \mathbb{Z}$ . Thus, we consider the following BNF for rational numbers

`<rat> ::= (cons <integer> <integer>)`

Define the following operations:

- (a) `(rat? q)`, which recognizes if  $q \in \langle \text{rat} \rangle$ .
  - (b) `(qsum q r)`, `(qdif q r)`, `(qmul q r)`, and `(qdiv q r)`, which take as inputs two values  $q, r \in \langle \text{rat} \rangle$  and compute their rational sum, difference, product, and division.
  - (c) `(rat-eq? q1 q2)`, which returns `#t` if  $q$  and  $r$  represent the same rational number, and `#f` otherwise.
  - (d) `(simplify q)`, which returns  $r \in \langle \text{rat} \rangle$  such that  $q$  and  $r$  represent the same rational number, and  $r = (\text{cons } a \ b)$  where  $a, b$  are relatively prime. To define this function, you can use the predefined function `(gcd u v)` which returns the greatest common divisor of two integers  $u, v \in \mathbb{Z}$ .
8. A complex number is a number whose value coincides with  $a + i \cdot b$  where  $a, b$  are floating-point numbers and  $i \in \mathbb{C}$  is the imaginary unit that satisfies the equation  $i^2 = -1$ . Suppose we choose to represent such a complex number as a pair `(cons a b)`. Thus, we consider the following BNF for complex numbers

$\langle \text{cplx} \rangle ::= (\text{cons } \langle \text{real} \rangle \langle \text{real} \rangle)$

Define the following operations:

- (a) `(cplx? c)`, which recognizes if  $c \in \langle \text{cplx} \rangle$ .
  - (b) `(abs-value c)`, which computes the absolute value of  $c \in \langle \text{cplx} \rangle$ . Remember that the absolute value of a complex number  $z = a + i \cdot b$  is  $|z| = \sqrt{a^2 + b^2}$ .
  - (c) `(rdiv c r)` which takes as inputs  $c \in \langle \text{cplx} \rangle$  and  $r \in \langle \text{real} \rangle$ ,  $r \neq 0$ , and returns the value from  $\langle \text{cplx} \rangle$  corresponding to the division of  $c$  by  $r$ .
  - (d) `(csum q r)`, `(cdif q r)`, `(cmul q r)`, and `(cdiv q r)`, which take as inputs two values  $q, r \in \langle \text{cplx} \rangle$  and compute their complex sum, difference, product, and division.
9. Consider sets represented by lists without repeating elements, and let  $A, B$  be two such sets. Define the following operations:
- (a) `(union A B)`, which computes the set union of  $A$  and  $B$ .
  - (b) `(difference A B)`, which computes the set difference of  $A$  and  $B$ .

You can use the predefined function `(member v lst)` which returns `#f` if  $v$  is not equal to any element of list `lst`, and true otherwise.