

## Applications of Disjoint Set Data Structures

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We describe here an application that can be solved easily using a disjoint-set data structure: the computation of a minimum-weight spanning tree of a weighted graph.

**Weighted graphs**

A **weighted graph** is a finite set of nodes connected by edges which have positive real numbers as weights. For example, the following is a weighted graph with 5 nodes and 6 edges: We will assume that

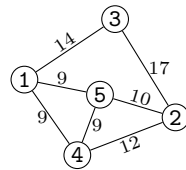


Figure 1: A weighed graph which is connected

- the nodes of a graph with  $n$  nodes are labeled with numbers from 1 to  $n$ .
- a text file which stores the representation of a weighted graph in the following way:
  - The first line contains the value of  $n$  (an integer)
  - The following lines contain 3 numbers separated by whitespace:  
 $i \quad j \quad w$   
 to indicate that the graph has an edge from node  $i$  to node  $j$  with weight  $w$ .

We assume that the edges are enumerated in increasing order of weight. For example, the weighted graph from Fig. 1 can be stored and read from a text file with the following content:

5  
 1 4 9  
 1 5 9  
 4 5 9  
 2 5 10  
 2 4 12  
 2 3 17

### Minimum weight spanning trees

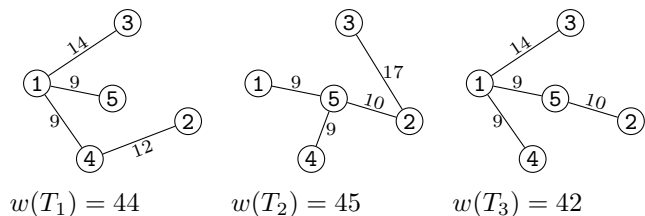
A graph is **connected** if there is a path between every two nodes in the graph. For example the weighted graph from Fig. 1 is connected.

A **spanning tree** of a weighted and connected graph  $G$  is a set  $T$  of edges of  $G$  such that

1. Every node of  $G$  is an endpoint of an edge in  $T$
2.  $T$  has no loops.

The **weight**  $w(T)$  of  $T$  is the sum of weights of edges in  $T$ .

For example, the following are spanning trees of the graph in Fig. 1:



A **minimum weight spanning tree** (or **MWST**) of  $G$  is a spanning tree of  $G$  whose weight has minimum possible value. For example,  $T_3$  is a MWST of the graph from Fig. 1.

A MWST of a connected and weighted graph  $G$  with  $n$  nodes can be found with Kruskal algorithm:

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Start with the initial partition  $S = \{\{1\}, \{2\}, \dots, \{n\}\}$ ,  $T = \emptyset$  and  $W = 0$ 
for each edge  $(i, j, w)$  of  $G$ , in increasing order of weights do
  if  $i, j$  are not in the same component of  $S$ 
    add  $(i, j, w)$  to  $T$ 
    Union( $i, j$ )
     $W = W + w$ 
  end if
end for
return  $T, W$ 

```

## Homework 1

Implement a program that reads from a text file `graph.txt` the representation of a connected weighted graph  $G$  and computes a MWST of  $G$ . The program will print the weight and the list of edges of the MWST.

**Suggestion:** implement a disjoint set-data structure using the explanations from Lecture 2, and use it to implement Kruskal algorithm.

## Homework 2

Indicate other possible applications of disjoint-set data structures.