

# Course 8

# Equivalence and Minimization of DFAs



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# Applications of interest

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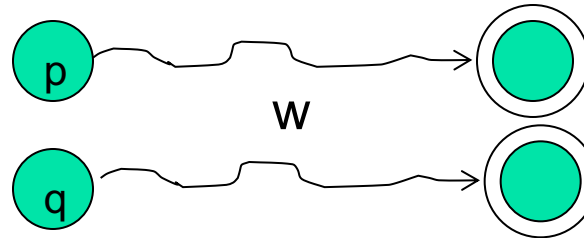
- Comparing two DFAs:
  - $L(\text{DFA}_1) == L(\text{DFA}_2)$ ?
  
- How to minimize a DFA?
  1. Remove unreachable states
  2. Identify & condense equivalent states into one

# When to call two states in a DFA “equivalent”?

Past doesn't matter - only future does!

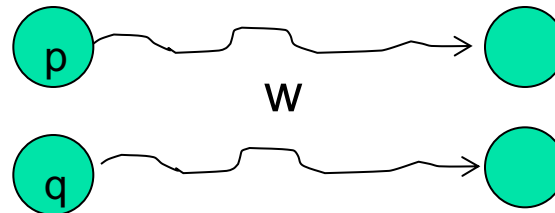
Two states  $p$  and  $q$  are said to be *equivalent* iff:

- i) Any string  $w$  accepted by starting at  $p$  is also accepted by starting at  $q$ ;



AND

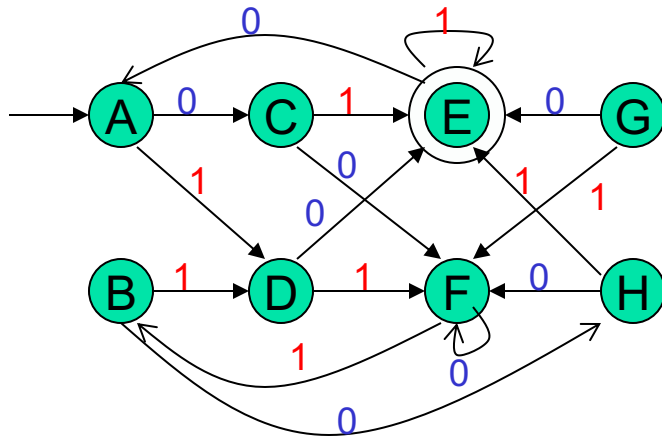
- i) Any string  $w$  rejected by starting at  $p$  is also rejected by starting at  $q$ .



→  $p \equiv q$

# Computing equivalent states in a DFA

## Table Filling Algorithm



### Pass #0

1. Mark accepting states  $\neq$  non-accepting states

### Pass #1

1. Compare every pair of states
2. Distinguish by one symbol transition
3. Mark = or  $\neq$  or blank (i.e. can not distinguish)

### Pass #2

1. Compare every pair of states
2. Distinguish by up to two symbol transitions (until different or same or tbd)

....

(keep repeating until table complete) *How the table on the right was obtained?* **Table Filling Algorithm** 4

A	=							
B	=	=						
C	x	x	=					
D	x	x	x	=				
E	x	x	x	x	=			
F	x	x	x	x	x	=		
G	x	x	x	=	x	x	=	
H	x	x	=	x	x	x	x	=
	A	B	C	D	E	F	G	H

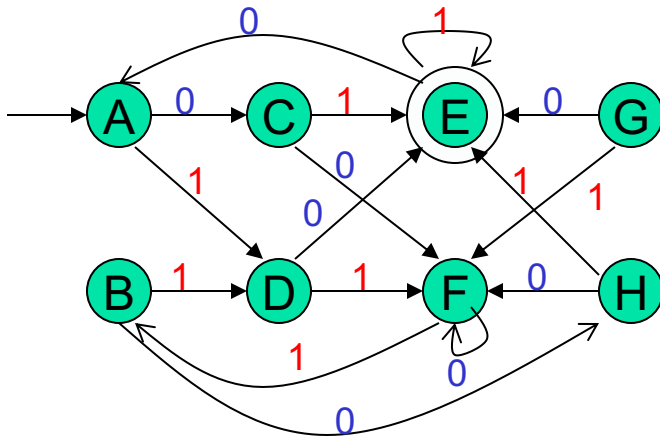


# Table Filling Algorithm

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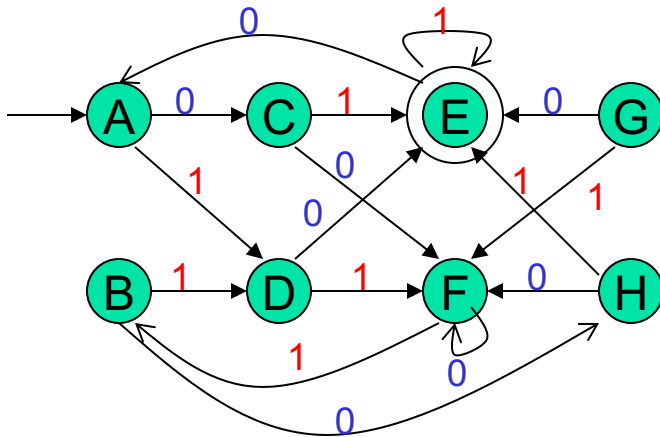
- Recursive discovery of distinguishable states in a DFA
  - **Base case:** If  $p$  is an accepting state and  $q$  is not accepting then the pair  $\{p, q\}$  is distinguishable.
  - **Induction:** Let  $p, q$  be states s.t. for some input symbol  $a$ ,  $r = \delta(p, a)$  and  $s = \delta(q, a)$  are known to be distinguishable. Then the pair  $\{p, q\}$  is distinguishable.

# Table Filling Algorithm - step by step



A	=							
B		=						
C			=					
D				=				
E					=			
F						=		
G							=	
H								=
	A	B	C	D	E	F	G	H

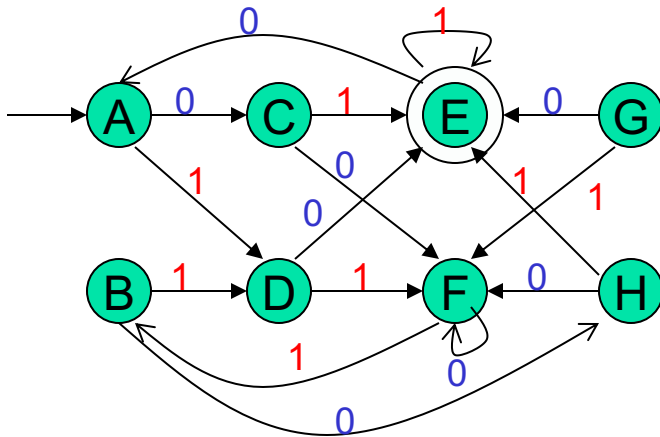
# Table Filling Algorithm - step by step



1. Mark X between accepting vs. non-accepting state

A	=							
B		=						
C			=					
D				=				
E	X	X	X	X	=			
F					X	=		
G					X		=	
H					X			=
	A	B	C	D	E	F	G	H

# Table Filling Algorithm - step by step



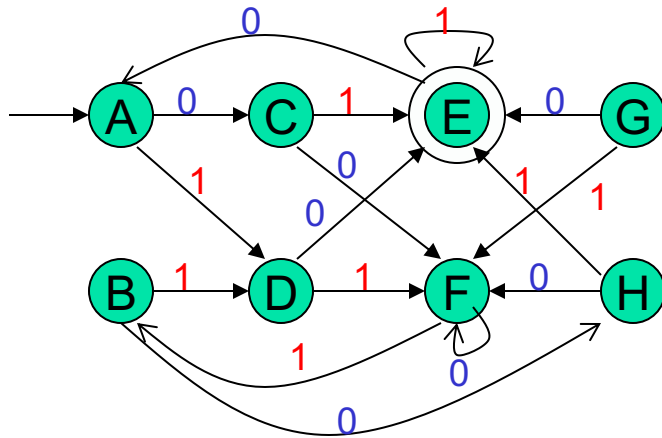
1. Mark X between accepting vs. non-accepting state
2. Look 1- hop away for distinguishing states or strings

A	=							
B		=						
C	X		=					
D	X			=				
E	X	X	X	X	=			
F					X	=		
G	X				X		=	
H	X				X			=
	A	B	C	D	E	F	G	H

↑



# Table Filling Algorithm - step by step

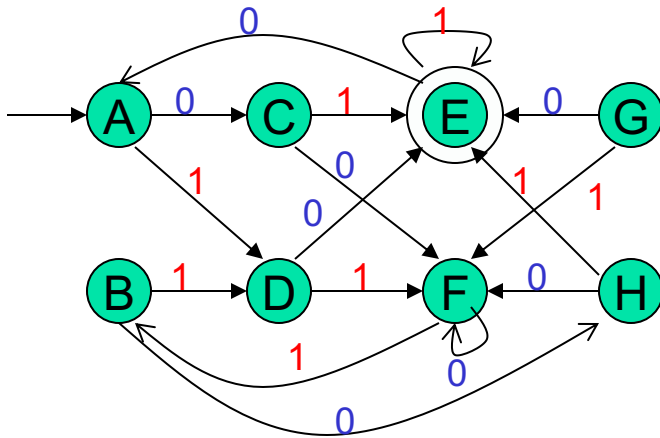


1. Mark X between accepting vs. non-accepting state
2. Look 1- hop away for distinguishing states or strings

A	=							
B		=						
C	X	X	=					
D	X	X		=				
E	X	X	X	X	=			
F					X	=		
G	X	X			X		=	
H	X	X			X			=
	A	B	C	D	E	F	G	H



# Table Filling Algorithm - step by step

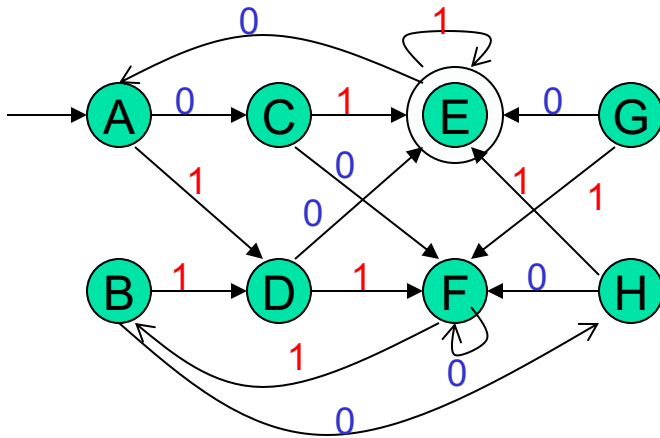


1. Mark X between accepting vs. non-accepting state
2. Look 1- hop away for distinguishing states or strings

A	=							
B		=						
C	X	X	=					
D	X	X	X	=				
E	X	X	X	X	=			
F			X		X	=		
G	X	X	X		X		=	
H	X	X	=		X			=
	A	B	C	D	E	F	G	H

↑

# Table Filling Algorithm - step by step

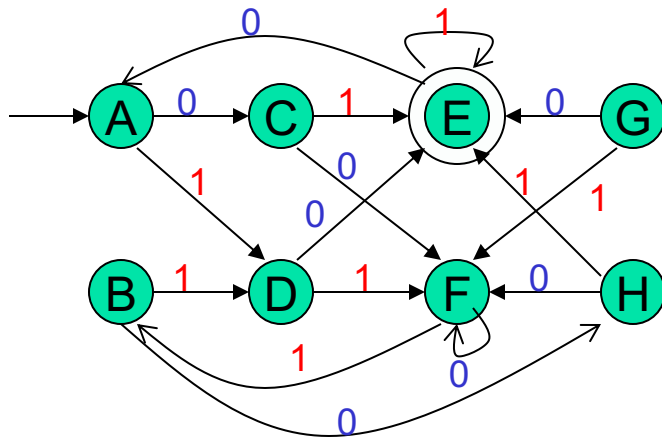


1. Mark X between accepting vs. non-accepting state
2. Look 1- hop away for distinguishing states or strings

A	=							
B		=						
C	X	X	=					
D	X	X	X	=				
E	X	X	X	X	=			
F			X	X	X	=		
G	X	X	X	=	X		=	
H	X	X	=	X	X			=
	A	B	C	D	E	F	G	H



# Table Filling Algorithm - step by step

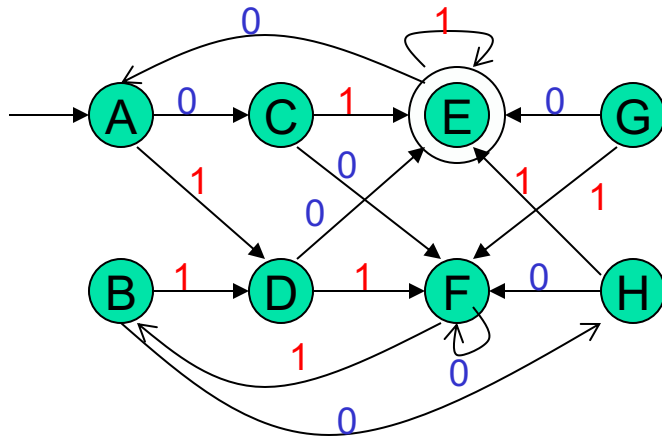


1. Mark X between accepting vs. non-accepting state
2. Look 1- hop away for distinguishing states or strings

A	=							
B		=						
C	X	X	=					
D	X	X	X	=				
E	X	X	X	X	=			
F			X	X	X	=		
G	X	X	X	=	X	X	=	
H	X	X	=	X	X	X		=
	A	B	C	D	E	F	G	H



# Table Filling Algorithm - step by step

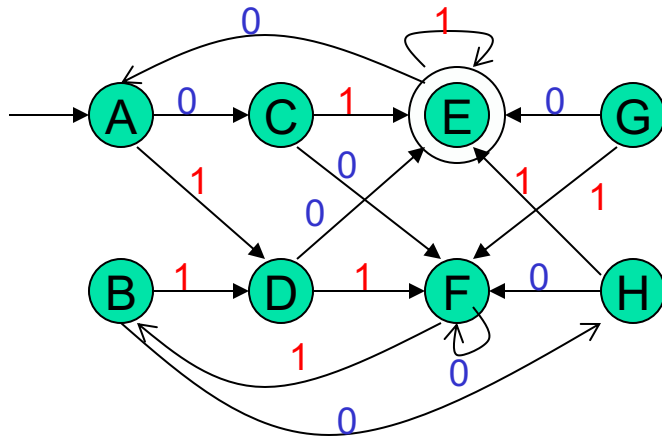


1. Mark X between accepting vs. non-accepting state
2. Look 1- hop away for distinguishing states or strings

A	=							
B		=						
C	X	X	=					
D	X	X	X	=				
E	X	X	X	X	=			
F			X	X	X	=		
G	X	X	X	=	X	X	=	
H	X	X	=	X	X	X	X	=
	A	B	C	D	E	F	G	H



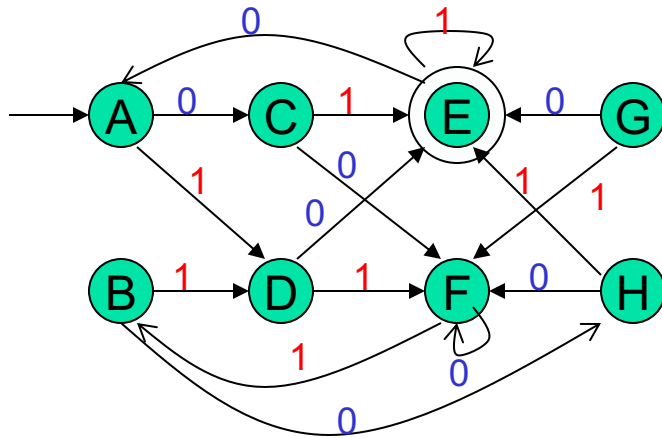
# Table Filling Algorithm - step by step



A	=							
B	=	=						
C	X	X	=					
D	X	X	X	=				
E	X	X	X	X	=			
F	X	X	X	X	X	=		
G	X	X	X	=	X	X	=	
H	X	X	=	X	X	X	X	=
	A	B	C	D	E	F	G	H

1. Mark **X** between accepting vs. non-accepting state
2. Pass 1:  
Look 1- hop away for distinguishing states or strings
3. Pass 2:  
Look 1-hop away again for distinguishing states or strings  
continue....

# Table Filling Algorithm - step by step



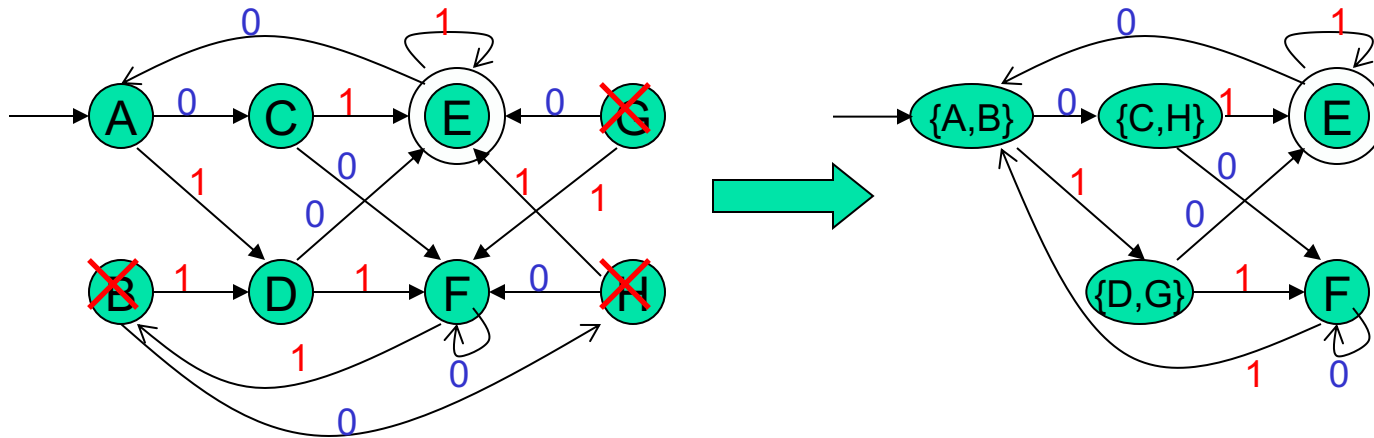
A	=							
B	=	=						
C	X	X	=					
D	X	X	X	=				
E	X	X	X	X	=			
F	X	X	X	X	X	=		
G	X	X	X	=	X	X	=	
H	X	X	=	X	X	X	X	=
	A	B	C	D	E	F	G	H

1. Mark X between accepting vs. non-accepting state
2. Pass 1:  
Look 1-hop away for distinguishing states or strings
3. Pass 2:  
Look 1-hop away again for distinguishing states or strings  
continue....

## Equivalences:

- A=B
- C=H
- D=G

# Table Filling Algorithm - step by step



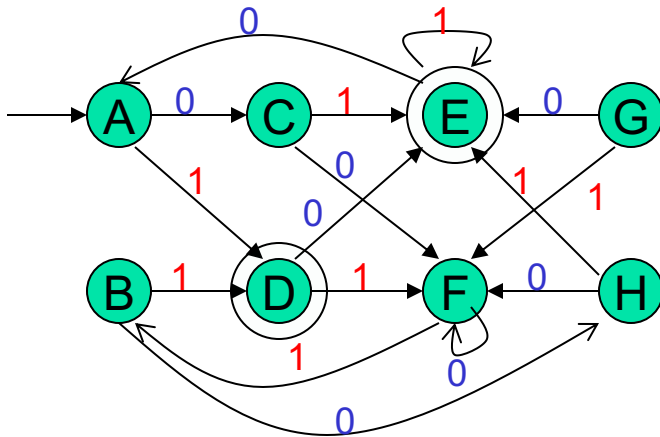
Retrain only one copy for each equivalence set of states

Equivalences:

- A=B
- C=H
- D=G



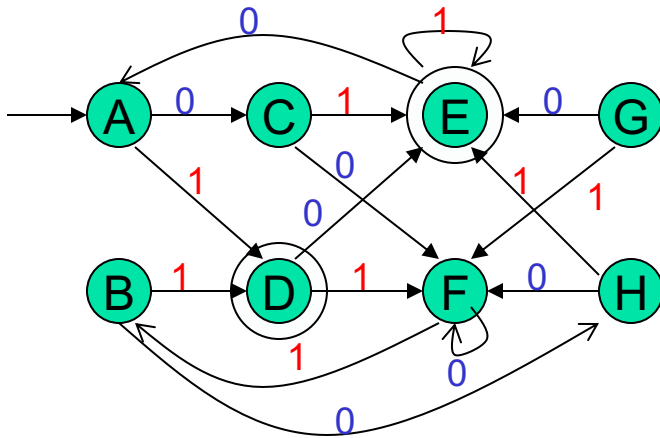
# Table Filling Algorithm – special case



A	=							
B		=						
C			=					
D				=				
E				?	=			
F						=		
G							=	
H								=
	A	B	C	D	E	F	G	H

Q) What happens if the input DFA has more than one final state?  
 Can all final states initially be treated as equivalent to one another?

# DFA Minimization by state equivalence method



0-Equivalence:  $\{A, B, C, F, G, H\}$   $\{D, E\}$   
 1-Equivalence  $\{A, B, C, H\}$   $\{F\}$   $\{G\}$   $\{D\}$   $\{E\}$

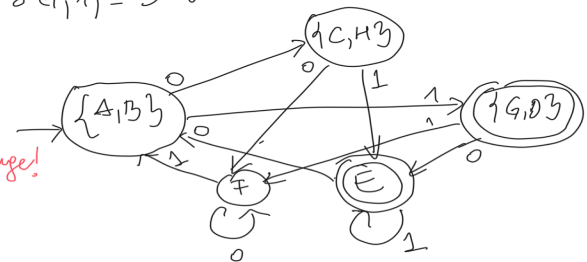
$\delta(A, 0) = C$   
 $\delta(B, 0) = A$   
 $\delta(A, 1) = D$   
 $\delta(B, 1) = D$   
 $\delta(C, 0) = E$   
 $\delta(C, 1) = F$   
 $\delta(D, 0) = E$   
 $\delta(D, 1) = F$   
 $\delta(E, 0) = G$   
 $\delta(E, 1) = E$   
 $\delta(F, 0) = H$   
 $\delta(F, 1) = A$   
 $\delta(G, 0) = E$   
 $\delta(G, 1) = F$   
 $\delta(H, 0) = F$   
 $\delta(H, 1) = E$

$\delta(H, 0) = F = \delta(C, 0)$   
 $\delta(H, 1) = E = \delta(C, 1)$   
 $\delta(D, 0) = E$   
 $\delta(E, 0) = A$

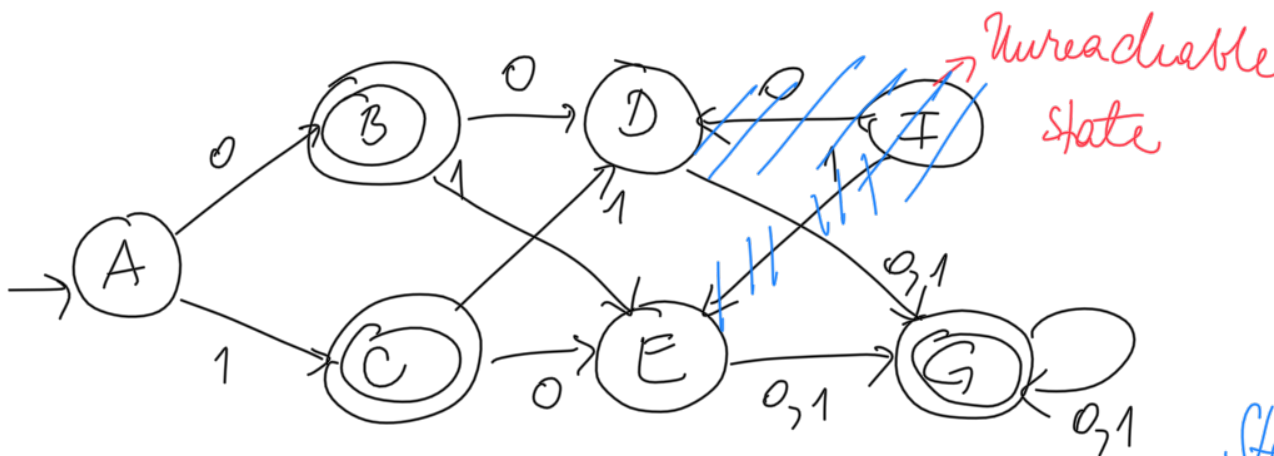
2-Equivalence  $\{A, B\}$   $\{C, H\}$   $\{F\}$   $\{G, D\}$   $\{E\}$

$\delta(A, 0) = C$   
 $\delta(H, 0) = F$   
 $\delta(C, 0) = E$   
 $\delta(F, 0) = F$   
 $\delta(G, 0) = E$   
 $\delta(H, 1) = E$   
 $\delta(C, 1) = F$   
 $\delta(F, 1) = B$   
 $\delta(D, 1) = F$   
 $\delta(G, 1) = F$

3-Equivalence  
 sets don't change!



# DFA Minimization with unreachable states



Step 1. Eliminate the unreachable state

Step 2. Proceed with state equivalence or table filling methods.

Putting it all together ...

## How to minimize a DFA?

- Goal: Minimize the number of states in a DFA
- Algorithm:
  - 1. Eliminate states unreachable from the start state
  - 2. Identify and remove equivalent states
  - 3. Output the resultant DFA

Depth-first traversal from the start state

Table filling algorithm



# Summary

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- Simplification of DFAs
  - How to remove unreachable states?
  - How to identify and collapse equivalent states?
  - How to minimize a DFA?